A circular parameterization for multi-sided patches

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Outline

Motivation Preliminaries Related work Circular domain Inverse problem Examples Algorithm Application Overlap-GB patch Conclusions







Why genuine multi-sided patches?



Parameterization



Domain evolution



Parallel tangents - periodic boundaries?

- Handled as multiple ribbons (CD)
- Handled as a single ribbon (GBS)
- Harmonic parameterization \Rightarrow discrete solution
- High computation cost, not suitable for some patch types



Parallel tangents - periodic boundaries?

Handled as multiple ribbons (CD) Handled as a single ribbon (GBS) Harmonic parameterization \Rightarrow discrete s High computation cost survable for some

Preliminaries

Properties I – Height parameter basics [linear map]

- 1. h = 0 on the base side.
- 2. *h* is continuous and varies monotonically.
- 3. *h* changes uniformly from 0 to 1 on the sides adjacent to the base side.



Properties II – Limited [interconnected map]

4. $h \leq 1$ everywhere inside the domain.



Properties III – Full [Wachspress map]

5. h = 1 on all distant sides. [full mapping]



Properties IV – Constrained [interconnected map]

6. $h'_{i-1} = -h'_{i+1}$ on the *i*th side. [constrained mapping]



All properties?

- Constrained Wachspress coordinates
- 1D multi-sided patch
- Singular blending function



All properties?



Circular domain

Inverse map



▶ Unit circle with equal arcs ▶ Base: $\left[-\frac{\pi}{n}, \frac{\pi}{n}\right]$ ▶ $\varphi = \frac{(2h+1)\pi}{n}$ ▶ $\theta = h\pi \implies \theta : 0 \to \pi$ ▶ $\psi = \theta - \varphi$ ▶ $\mathbf{0} = \left(\frac{\sin\theta}{\sin\psi}, 0\right)$ ▶ $r = \left|\frac{\sin\varphi}{\sin\psi}\right|$





Constrained property & corner parameterization





Algorithm

• Line at
$$\hat{h} = \frac{1}{n-2} \Rightarrow \hat{u} = \cos \frac{\pi}{n-2}$$

• Same circle for $h = 0$ and $h = 1$
• Idea: bisection search

 $\begin{array}{l|l} \text{if } u > \hat{u} \text{ then} \\ | & \text{return bisection}(\Delta, 0, \hat{h} - \varepsilon) \\ \text{if } u < \hat{u} \text{ then} \\ | & \text{return bisection}(\Delta, \hat{h} + \varepsilon, 1) \\ \text{return } \hat{h} \end{array}$

 $\Delta(h) = \|\mathbf{p} - \mathbf{O}(h)\| - r(h)$



Algorithm

if

if

▶ Line at
$$\hat{h} = \frac{1}{n-2} \Rightarrow \hat{u} = \cos \frac{\pi}{n-2}$$

▶ Same circle for $h = 0$ and $h = 1$
▶ Idea: bisection search
if $u > \hat{u}$ then
| return bisection($\Delta, 0, \hat{h} - \varepsilon$)
if $u < \hat{u}$ then
| return bisection($\Delta, \hat{h} + \varepsilon, 1$)
return \hat{h}

 $\Delta(h) = \|\mathbf{p} - \mathbf{O}(h)\| - r(h)$



Application

Overlap–GB patch

Corner-based variation of the GB patch

Needs a full, constrained parameterization

$$S = \sum_{i=1}^{n} \sum_{j=0}^{\lfloor d/2 \rfloor} \sum_{k=0}^{\lfloor d/2 \rfloor} \mathbf{P}_{ijk} B_j^d(h_{i+1}) B_k^d(h_i) + \mathbf{P}_0 B_0,$$



Conclusions

- Circle as multi-sided domain
- Height parameterization
 - Circular arcs
 - Full
 - Constrained
 - Efficient
- Overlap–GB patch
- Suitable for periodic boundaries
- ► *G*² cap for subdivision surfaces?



Conclusions

- Circle as multi-sided domain
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Any questions?

