**Tessellation of Zheng–Ball patches**

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### Abstract

Some of the earliest multi-sided surface formulations depend on a parameterization defined implicitly by constraints. This presents a difficulty in visualizing the patches, as the parametric domain is not trivial to triangulate. We propose a general method that can be used with any number of sides.

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1. **Introduction**

In recent years, n-sided surface representations have received renewed attention. These formulations can describe non-four-sided free-form surfaces without the inaccuracies of the usual trimming approach. An example of such patches was proposed by Sabin, where 2n + 1 control points were used to define three- and five-sided surfaces that could connect to quadratic four-sided patches with $G^1$ continuity. The same idea was extended later to six-sided surfaces and cubic boundaries, and then to a general degree.

These surfacing methods depend on n scalar parameters defined implicitly by constraints, representing a kind of distance from the boundaries. Consequently, the patch equation is a mapping from $R^n$ to $R^2$, and to visualize the surface, we first need a tessellation of the n-dimensional parameter values. It is not evident how to do this, and this paper explores different options.

2. **Preliminaries**

Parameters $u = \{u_i\}^n_i \in [0, 1]^n$ form a 2-dimensional surface in $R^n$ space. The side of the domain where $u_i = 0$ is mapped to the $i^{th}$ boundary of the patch. At these points, the following constraints apply:

$$u_{i-1} + u_{i+1} = 1, \quad u_j = 1, \quad j \neq i-1, i, i+1$$ (1)

$$\sum_j u_j = c_n u_{i-1} u_{i+1}$$ (2)

$$u_{i+1} u'_{i-1} = u_{i-1} u'_{i+1}, \quad \text{when } n > 3$$ (3)

where $c_n$ is a constant depending only on the number of sides, and derivatives are with respect to $u_i$. The patch equation using these parameters can be found in the Appendix.

3. **Tessellation**

In the following we will review the parameterizations of three-, five- and six-sided patches, and propose tessellation algorithms of increasing complexity to handle the arising difficulties.

3.1. **Three-sided patches**

The three-sided case is the simplest. Using the normalization equation

$$u_1 + u_2 + u_3 - 2u_1 u_2 u_3 = 1, \quad (5)$$

we can choose $c_3 = 2, u'_{i-1} = -u_{i-1}^2$ and $u'_{i+1} = -u_{i+1}^2$ to satisfy Eqs. (1)–(3). By Eq. (2) we also know that corners of the domain have parameters $(1, 0, 0), (0, 1, 0), (0, 0, 1)$ – generally sequences of 1s with two consecutive 0s at the indices associated with the two boundaries at the corner. We can create a linear triangulation connecting these three vertices, see Fig. 1a.

Now we can take two coordinates, e.g. $u_1$ and $u_2$, and express the third by them using Eq. (5):

$$\hat{u}_3 = \frac{1 - u_1 - u_2}{1 - 2u_1 u_2}. \quad (6)$$

The $(u_1, u_2, \hat{u}_3)$ points are on the domain surface, see Figure 1b. The difference from the linear (barycentric) parameters can be seen clearly on Fig. 2, where the domain is depicted from the side. Note that the triangulation will not be symmetric, since we arbitrarily chose $\hat{u}_3$ to express from the other two, but this is not important at sufficiently high resolutions.
3.2. Five- and six-sided patches

For five-sided patches, we have the normalization equations:

$$u_j = 1 - u_{j+2}u_{j+3}, \quad j = 1 \ldots 5 \quad (7)$$

using cyclic indexing. At first sight, this defines 5 equations for the 5 parameters, but actually only 3 of the equations are independent, so we still get a two-dimensional surface in $\mathbb{R}^5$. Setting $c_5 = 1$, $u'_{i-1} = u_{i-1}u_{i-1}$ and $u'_{i+1} = u_{i-1}u_{i-1}$, it is easy to see that Eqs. (1)–(4) hold.

An $n$-sided polygon can be tessellated by first dividing it into $n$ triangles, see Fig. 3. For this, we need a center position. It is natural to require that at the middle of the surface all parameters should be equal. Simple algebra shows that $(\varphi, \varphi, \varphi, \varphi, \varphi)$ is actually on the domain surface, where $\varphi = (\sqrt{5} - 1)/2$ is the golden ratio.

We can create a linear triangulation of each sub-triangle, as in the previous section. Choosing $u_1$ and $u_2$ as base, the rest can be computed as

$$\hat{u}_3 = \frac{1 - u_1}{1 - u_1u_2}, \quad \hat{u}_4 = 1 - u_1u_2, \quad \hat{u}_5 = \frac{1 - u_2}{1 - u_1u_2}, \quad (8)$$

but then the sub-triangle associated with the $u_4 = 0$ boundary will be degenerate, as on that side $u_1 = u_2 = 1$.

It is better to “project” the linear triangulation onto the domain surface by minimizing the normalization energy

$$E_5(u) = \sum_{j=1}^{5} (u_j + u_{j+2}u_{j+3} - 1)^2, \quad (9)$$

which is just the sum of squared normalization errors based on Eq. (7). This can be minimized using a simple derivative-free optimization algorithm, such as Powell’s method. Figure 4 shows constant parameter lines before and after the projection. While these show various breaks, remember that this represents only one coordinate of a continuous 5-dimensional surface.

Instead of starting from a linear approximation, we can also make use of generalized barycentric coordinates. First we tessellate a regular $n$-sided polygon (as in Fig. 3), then
map the vertices to (e.g.) Wachspress coordinates \( w = \{ w_j \} \). We can use the distance parameter \( s = 1 - w_i - w_{i-1} \) as our first approximation, since this also satisfies Eqs. (1)–(2). Finally the vertices are projected onto the domain as above, so that Eqs. (3)–(4) will also hold. Figure 5 shows constant parameter lines for this method. Note that the parameters in Figures 4 and 5 define the same surface.

Generalized barycentric coordinates can also be used with the direct method of Eq. (8). The \( u_4 = 0 \) side will still be degenerate (we can use \( u_4 = u_3 \) there) but other parts of the sub-triangle behave normally.

In the case of six-sided patches, there is no rational formula for the parameters.\(^7\) The normalization equations (only four of which are independent) are:\(^7\)

\[
\begin{align*}
  u_{i+1}^2 (1 - u_{i-1} u_i) (1 - 2u_{i-1} u_i) + \\
  u_{i+1} (2u_{i-1} - 3u_{i-1}^2 u_i + u_{i-1} u_i^2) + u_{i+1}^2 &= 1. \\
  (10)
\end{align*}
\]

Setting \( c_6 = 1 \) with

\[
\begin{align*}
  u_i'_{i-2} &= \frac{1}{2} u_{i-1} u_{i+1} - u_{i+1}, \\
  (11) \\
  u_i'_{i-1} &= \frac{3}{2} u_{i-1} u_{i+1}, \\
  u_i'_{i+1} &= \frac{3}{2} u_{i-1} u_{i+1}, \\
  (12) \\
  u_i'_{i+2} &= \frac{1}{2} u_{i-1} u_{i+1} - u_{i-1}, \\
  (13) \\
  u_i'_{i+3} &= -\frac{1}{2} u_{i-1} u_{i+1}, \\
  (14)
\end{align*}
\]

all the requirements are satisfied. Since the direct method would involve solving a non-linear equation, it is simpler to use projection here, as well. The center point in this case is \((\psi, \psi, \psi, \psi, \psi)\) with \( \psi = \frac{1}{\sqrt{2}} \).

4. Results

Figure 6 shows a 5-sided patch triangulated with the projection method, starting from both the linear and Wachspress-based approximations, with 10 triangles on each boundary. Generalized Bézier patches\(^9\) share the same control struc-
ture, see Fig. 7. A comparison is shown on Fig. 8, with a linear-based dense triangulation (50 triangles per boundary). The patches are generated from the same control points. The isophote lines are very similar, and match at the boundaries (as required).

Conclusion
We have shown different methods for the tessellation of Zheng–Ball patches. Of these, the only one that scales to arbitrary number of sides is the projection method. As future work, gradient descent could be used instead of Powell’s method, which would result in increased symmetry. Parameterizations for more sides should also be investigated, probably using biharmonic surfaces.

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Appendix A: The Zheng–Ball patch equation
All control points are given indices based on their position in the control structure. An index $\lambda = \{\lambda_j\}_1^n$ is composed of $n$ natural numbers, with $\lambda_j$ signifying how far removed is the control point from the $i$th side. For example, a control point on side $i$ has $\lambda_i = 0$; the central control point in Fig. 7a has the index $(2,2,2,2,2,2)$. Generally, in a patch of $n > 3$ sides with degree-$d$ boundaries, when the $(i-1)^{st}$ and $i^{th}$ sides are closest to a control point, the remaining indices are computed as

$\lambda_{i-2} = d - \lambda_i$, $\lambda_{i+1} = d - \lambda_{i-1}$, $\lambda_j = d - \min(\lambda_{i-1}, \lambda_{i+1})$, for $j \notin \{i - 2, i - 1, i, i + 1\}$.

The patch is evaluated at parameter $u$ as

$S(u) = \sum_{\lambda} P_{\lambda} B_{\lambda}(u)$,

where $P_{\lambda}$ denotes the control points, and $B_{\lambda}(u)$ is a suitable blending function.

For control points not on the boundary (i.e., $\forall j : \lambda_j > 0$) we define the following blend:

$B_{\lambda}(u) = \binom{d}{\lambda_{i-1}} \binom{d}{\lambda_i} \prod_{j=1}^{n} u_j^{\lambda_j}$,

where $\lambda_{i-1}$ and $\lambda_i$ are the smallest elements in $\lambda$. For control points on the $i^{th}$ boundary, we have $\lambda_{i-1} = k$, $\lambda_i = 0$, $\lambda_{i+1} = d - k$, and all other values in $\lambda$ are 1. In this case, the blend is defined as

$B_{\lambda}(u) = \binom{d}{k} \prod_{j \neq i}^{n} u_j^{\lambda_j} \prod_{j=1}^{n} u_j^{\lambda_j} \left(1 - d c_n \prod_{j=1}^{n} u_j\right)$.
This latter equation is a bit different for triangular patches:

\[ B^*_\lambda(u) = \binom{d}{k} \frac{u_{d-1}^k u_{d+1}^{d-k}}{(1 - \lambda f_{\lambda}(u))} \]

where

\[ f_{\lambda}(u) = \begin{cases} 
  d - k + (2k - d)u_{i+1} & \text{when } d - k \leq k, \\
  k + (d - 2k)u_{i-1} & \text{when } d - k > k.
\end{cases} \]

Unfortunately \( B^*_\lambda(u) = \sum_{\lambda} B^*_\lambda(u) \) does not equal to 1, which would be needed for the surface to be an affine combination of the control points, so we distribute the weight deficiency \((1 - B^*_\lambda(u))\) between the blending functions associated with control points not on the boundary:

\[ B^*_\lambda(u) = \begin{cases} 
  B^*_\lambda(u) & \text{when } \exists j : \lambda_j = 0, \\
  B^*_\lambda(u) + \frac{1 - B^*_\lambda(u)}{T_{n,d}} & \text{otherwise},
\end{cases} \]

where

\[ T_{n,d} = \begin{cases} 
  \frac{n(d-2)d}{2} + 1 & \text{when } d \text{ is even}, \\
  \frac{n(d-1)^2}{4} & \text{when } d \text{ is odd}
\end{cases} \]

is the number of control points not on the boundary.

References